# NAG Fortran Library Routine Document

# H03ADF

Note: before using this routine, please read the Users' Note for your implementation to check the interpretation of *bold italicised* terms and other implementation-dependent details.

## 1 Purpose

H03ADF finds the shortest path through a directed or undirected acyclic network using Dijkstra's algorithm.

## 2 Specification

```
SUBROUTINE H03ADF(N, NS, NE, DIRECT, NNZ, D, IROW, ICOL, SPLEN, PATH,1IWORK, WORK, IFAIL)INTEGERN, NS, NE, NNZ, IROW(NNZ), ICOL(NNZ), PATH(N),1IWORK(3*N+1), IFAILrealD(NNZ), SPLEN, WORK(2*N)LOGICALDIRECT
```

## **3** Description

This routine attempts to find the shortest path through a **directed** or **undirected acyclic** network, which consists of a set of points called **vertices** and a set of curves called **arcs** that connect certain pairs of distinct vertices. An acyclic network is one in which there are no paths connecting a vertex to itself. An arc whose origin vertex is i and whose destination vertex is j can be written as  $i \rightarrow j$ . In an undirected network the arcs  $i \rightarrow j$  and  $j \rightarrow i$  are equivalent (i.e.,  $i \leftrightarrow j$ ), whereas in a directed network they are different. Note that the shortest path may not be unique and in some cases may not even exist (e.g., if the network is disconnected).

The network is assumed to consist of n vertices which are labelled by the integers 1, 2, ..., n. The lengths of the arcs between the vertices are defined by the n by n **distance matrix** D, in which the element  $d_{ij}$  gives the length of the arc  $i \rightarrow j$ ;  $d_{ij} = 0$  if there is no arc connecting vertices i and j (as is the case for an acyclic network when i = j). Thus the matrix D is usually **sparse**. For example, if n = 4 and the network is directed, then

$$\mathbf{D} = \begin{pmatrix} 0 & d_{12} & d_{13} & d_{14} \\ d_{21} & 0 & d_{23} & d_{24} \\ d_{31} & d_{32} & 0 & d_{34} \\ d_{41} & d_{42} & d_{43} & 0 \end{pmatrix}$$

If the network is undirected, D is symmetric since  $d_{ij} = d_{ji}$  (i.e., the length of the arc  $i \to j \equiv$  the length of the arc  $j \to i$ ).

The method used by H03ADF is described in detail in Section 8.

## 4 References

Dijkstra E W (1959) A note on two problems in connection with graphs Numer. Math. 1 269-271

## 5 Parameters

1: N – INTEGER

On entry: n, the number of vertices. Constraint:  $N \ge 2$ . Input

2: NS – INTEGER

3: NE – INTEGER

On entry:  $n_s$  and  $n_e$ , the labels of the first and last vertices, respectively, between which the shortest path is sought.

Constraints:

 $\begin{array}{l} 1 \leq \text{NS} \leq \text{N}, \\ 1 \leq \text{NE} \leq \text{N}, \\ \text{NS} \neq \text{NE}. \end{array}$ 

#### 4: DIRECT – LOGICAL

On entry: indicates whether the network is directed or undirected as follows:

if DIRECT = .TRUE., the network is directed;

if DIRECT = .FALSE., the network is undirected.

5: NNZ – INTEGER

On entry: the number of non-zero elements in the distance matrix D.

Constraints:

if DIRECT = .TRUE.,  $1 \le NNZ \le N \times (N-1)$ ; if DIRECT = .FALSE.,  $1 \le NNZ \le N \times (N-1)/2$ .

6: D(NNZ) – *real* array

On entry: the non-zero elements of the distance matrix D, ordered by increasing row index and increasing column index within each row. More precisely, D(k) must contain the value of the non-zero element with indices (IROW(k), ICOL(k)); this is the length of the arc from the vertex with label IROW(k) to the vertex with label ICOL(k). Elements with the same row and column indices are not allowed. If DIRECT = .FALSE., then only those non-zero elements in the strict upper triangle of D need be supplied since  $d_{ij} = d_{ji}$ . (F11ZAF may be used to sort the elements of an arbitrarily ordered matrix into the required form. This is illustrated in Section 9.)

*Constraint*: D(k) > 0.0, for k = 1, 2, ..., NNZ.

- 7: IROW(NNZ) INTEGER array
- 8: ICOL(NNZ) INTEGER array

On entry: IROW(k) and ICOL(k) must contain the row and column indices, respectively, for the non-zero element stored in D(k).

#### Constraints:

IROW and ICOL must satisfy the following constraints (which may be imposed by a call to F11ZAF):

$$\begin{split} &\text{IROW}(k-1) < \text{IROW}(k) \text{, or} \\ &\text{IROW}(k-1) = \text{IROW}(k) \text{ and } \text{ICOL}(k-1) < \text{ICOL}(k) \text{, for } k = 2, 3, \dots, \text{NNZ.} \\ &\text{In addition, if } \text{DIRECT} = .\text{TRUE., } 1 \leq \text{IROW}(k) \leq \text{N}, \quad 1 \leq \text{ICOL}(k) \leq \text{N} \quad \text{and} \\ &\text{IROW}(k) \neq \text{ICOL}(k); \\ &\text{if } \text{DIRECT} = .\text{FALSE., } 1 \leq \text{IROW}(k) < \text{ICOL}(k) \leq \text{N}. \end{split}$$

On exit: IROW is used as internal workspace prior to being restored and hence is unchanged.

### 9: SPLEN – *real*

Output

On exit: the length of the shortest path between the specified vertices  $n_s$  and  $n_e$ .

Input

Input

Input

Input

Input/Output

Input

Input

[NP3546/20A]

10: PATH(N) – INTEGER array

On exit: contains details of the shortest path between the specified vertices  $n_s$  and  $n_e$ . More precisely, NS = PATH(1)  $\rightarrow$  PATH(2)  $\rightarrow \ldots \rightarrow$  PATH(p) = NE for some  $p \leq n$ . The remaining (n-p) elements are set to zero.

- 11: IWORK(3\*N+1) INTEGER array
- 12: WORK(2\*N) *real* array
- 13: IFAIL INTEGER

On entry: IFAIL must be set to 0, -1 or 1. Users who are unfamiliar with this parameter should refer to Chapter P01 for details.

On exit: IFAIL = 0 unless the routine detects an error (see Section 6).

For environments where it might be inappropriate to halt program execution when an error is detected, the value -1 or 1 is recommended. If the output of error messages is undesirable, then the value 1 is recommended. Otherwise, for users not familiar with this parameter the recommended value is 0. When the value -1 or 1 is used it is essential to test the value of IFAIL on exit.

### 6 Error Indicators and Warnings

If on entry IFAIL = 0 or -1, explanatory error messages are output on the current error message unit (as defined by X04AAF).

Errors or warnings detected by the routine:

IFAIL = 1

```
IFAIL = 2
```

#### IFAIL = 3

On entry, IROW(k) < 1 or IROW(k) > N or ICOL(k) < 1 or ICOL(k) > N or IROW(k) = ICOL(k) for some k when DIRECT = .TRUE..

### IFAIL = 4

On entry, IROW(k) < 1 or  $IROW(k) \ge ICOL(k)$  or ICOL(k) > N for some k when DIRECT = .FALSE.

### IFAIL = 5

 $D(k) \leq 0.0$  for some k.

## IFAIL = 6

On entry, IROW(k-1) > IROW(k) or IROW(k-1) = IROW(k) and ICOL(k-1) > ICOL(k) for some k.

Workspace

Workspace

Input/Output

Output

IFAIL = 7

On entry, IROW(k-1) = IROW(k) and ICOL(k-1) = ICOL(k) for some k.

#### IFAIL = 8

No connected network exists between vertices NS and NE.

## 7 Accuracy

The results are exact, except for the obvious rounding errors in summing the distances in the length of the shortest path.

### 8 **Further Comments**

This routine is based upon Dijkstra's algorithm (see Dijkstra (1959)), which attempts to find a path  $n_s \rightarrow n_e$  between two specified vertices  $n_s$  and  $n_e$  of shortest length  $d(n_s, n_e)$ .

The algorithm proceeds by assigning labels to each vertex, which may be **temporary** or **permanent**. A temporary label can be changed, whereas a permanent one cannot. For example, if vertex p has a permanent label (q, r), then r is the distance  $d(n_s, r)$  and q is the previous vertex on a shortest length  $n_s \rightarrow p$  path. If the label is temporary, then it has the same meaning but it refers only to the shortest  $n_s \rightarrow p$  path found so far. A shorter one may be found later, in which case the label may become permanent.

The algorithm consists of the following steps.

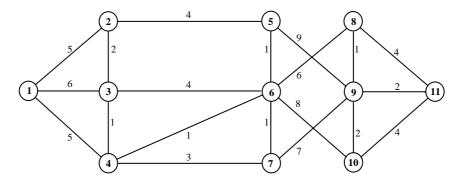
- 1. Assign the permanent label (-,0) to vertex  $n_s$  and temporary labels  $(-,\infty)$  to every other vertex. Set  $k = n_s$  and go to (2).
- 2. Consider each vertex y adjacent to vertex k with a temporary label in turn. Let the label at k be (p,q) and at y(r,s). If  $q + d_{ky} < s$ , then a new temporary label  $(k, q + d_{ky})$  is assigned to vertex y; otherwise no change is made in the label of y. When all vertices y with temporary labels adjacent to k have been considered, go to (3).
- 3. From the set of temporary labels, select the one with the smallest second component and declare that label to be permanent. The vertex it is attached to becomes the new vertex k. If  $k = n_e$  go to (4). Otherwise go to (2) unless no new vertex can be found (e.g., when the set of temporary labels is 'empty' but  $k \neq n_e$ , in which case no connected network exists between vertices  $n_s$  and  $n_e$ ).
- 4. To find the shortest path, let (y, z) denote the label of vertex  $n_e$ . The column label (z) gives  $d(n_s, n_e)$  while the row label (y) then links back to the previous vertex on a shortest length  $n_s \rightarrow n_e$  path. Go to vertex y. Suppose that the (permanent) label of vertex y is (w, x), then the next previous vertex is w on a shortest length  $n_s \rightarrow y$  path. This process continues until vertex  $n_s$  is reached. Hence the shortest path is

$$n_s \to \ldots \to w \to y \to n_e,$$

which has length  $d(n_s, n_e)$ .

## 9 Example

To find the shortest path between vertices 1 and 11 for the undirected network



### 9.1 Program Text

```
HO3ADF Example Program Text
*
      Mark 20 Revised. NAG Copyright 2001.
*
      .. Parameters ..
      INTEGER
                       NIN, NOUT
                       (NIN=5,NOUT=6)
      PARAMETER
      INTEGER
                       NMAX, NNZMAX
                        (NMAX=100,NNZMAX=1000)
     PARAMETER
                       DUP, ZERO
      CHARACTER
                       (DUP='F',ZERO='R')
     PARAMETER
      .. Local Scalars ..
*
     real
                       SPLEN
      INTEGER
                       IFAIL, J, LENC, N, NE, NNZ, NS
      LOGICAL
                       DIRECT
      .. Local Arrays ..
*
      real
                       D(NNZMAX), WORK(2*NMAX)
                        ICOL(NNZMAX), IROW(NNZMAX), IWORK(3*NMAX+1),
     INTEGER
                       PATH(NMAX)
      .. External Subroutines ...
*
     EXTERNAL
                       F11ZAF, H03ADF
      .. Executable Statements ..
     WRITE (NOUT, *) 'HO3ADF Example Program Results'
      Skip heading in data file
*
      READ (NIN, *)
      READ (NIN, *) N, NS, NE, NNZ, DIRECT
      IF (N.LE.NMAX .AND. NNZ.LE.NNZMAX) THEN
         Read D, IROW and ICOL from data file.
*
         READ (NIN,*) (D(J),IROW(J),ICOL(J),J=1,NNZ)
*
         Reorder the elements of D into the form required by HO3ADF.
*
*
         IFAIL = 0
         CALL F11ZAF(N,NNZ,D,IROW,ICOL,DUP,ZERO,IWORK,IWORK(N+2),IFAIL)
         Find the shortest path between vertices NS and NE.
*
         IFAIL = 0
         CALL H03ADF(N,NS,NE,DIRECT,NNZ,D,IROW,ICOL,SPLEN,PATH,IWORK,
     +
                     WORK, IFAIL)
*
         IF (IFAIL.EQ.0) THEN
*
            Print details of shortest path.
            DO 20 J = 0, N - 1
               IF (PATH(J+1).EQ.0) THEN
```

```
LENC = J
                 GO TO 40
              END IF
  20
           CONTINUE
           LENC = N
  40
           CONTINUE
           WRITE (NOUT,99999) 'Shortest path = ', (PATH(J),J=1,LENC)
           WRITE (NOUT,99998) 'Length of shortest path = ', SPLEN
        END IF
     END IF
     STOP
*
99999 FORMAT (/1X,A,10(I2,:' to '))
99998 FORMAT (/1X,A,G16.6)
     END
```

### 9.2 Program Data

H03ADF Example				Program Data									
11	1	11	20	F	:Valu	les	of	Ν,	NS,	NE,	NNZ	and	DIRECT
6.0	68												
1.0	8	89											
2.0	ç	9 11											
4.0	2												
1.0	3	3 4											
6.0	1	1 3											
4.0	З	36											
1.0	4	Į	6										
2.0	2	2	3										
3.0	4	ļ	7										
5.0	1	_	2										
7.0	6	5 1	.0										
1.0	5	56											
4.0	8	3 1	.1										
9.0	5	5	9										
1.0	6		7										
8.0	7	7	9										
4.0	10	) 1	.1										
2.0	ç	) 1	.0										
5.0	1	-	4		:End	of	D,	IRO	⊃W,	ICOL			

### 9.3 Program Results

HO3ADF Example Program Results

Shortest path = 1 to 4 to 6 to 8 to 9 to 11

Length of shortest path = 15.0000